

Memory Analysis and CPU-time Profiling in RMG-Java

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Outline

Memory Management in Java



Demo of Memory Analysis



CPU-time Profiling



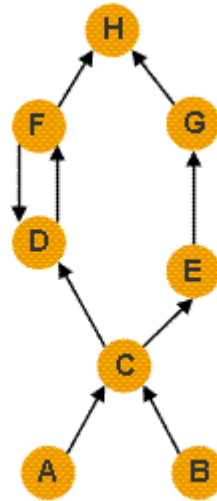
Possible Approaches

Memory Management in Java

- Programming languages like C/C++
 - Manually allocate/de-allocate memory
- Java
 - Automatically de-allocate
 - Garbage collector

Basic concepts for Garbage Collection

- Heap dump



- Shallow heap

Memory consumed by one object itself

- G.C. root

Any variables your program can access directly

- Local variables
- Class static variables

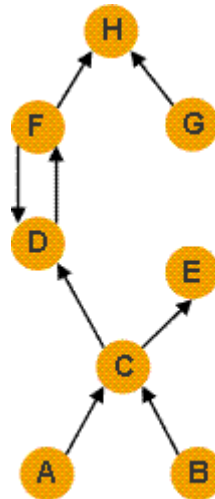
Basic concepts for Garbage Collection

- Live objects

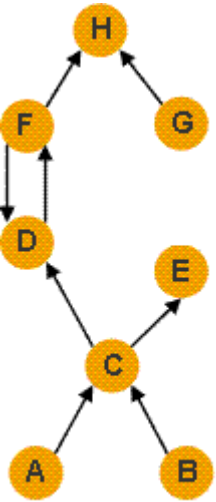
Can be reached from G.C. Root

- Retained set & heap

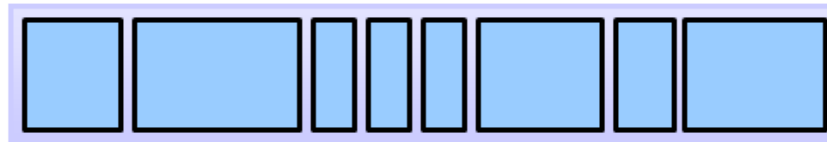
- Dominator



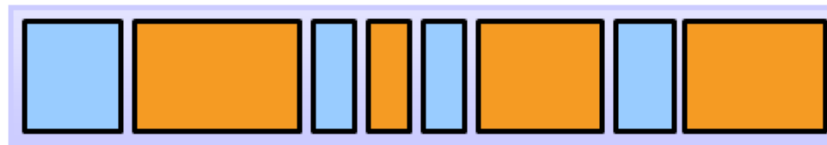
Mark and Sweep Garbage Collection






Marking



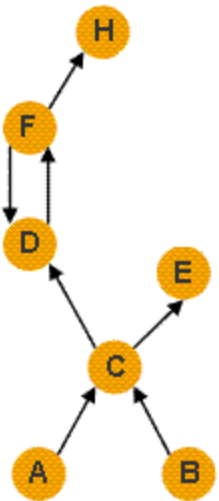
Before Marking



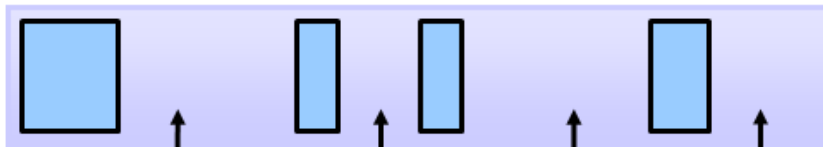
After Marking

-  A live object
-  Unreferenced Objects
-  Memory space

Mark and Sweep Garbage Collection



Normal Deletion

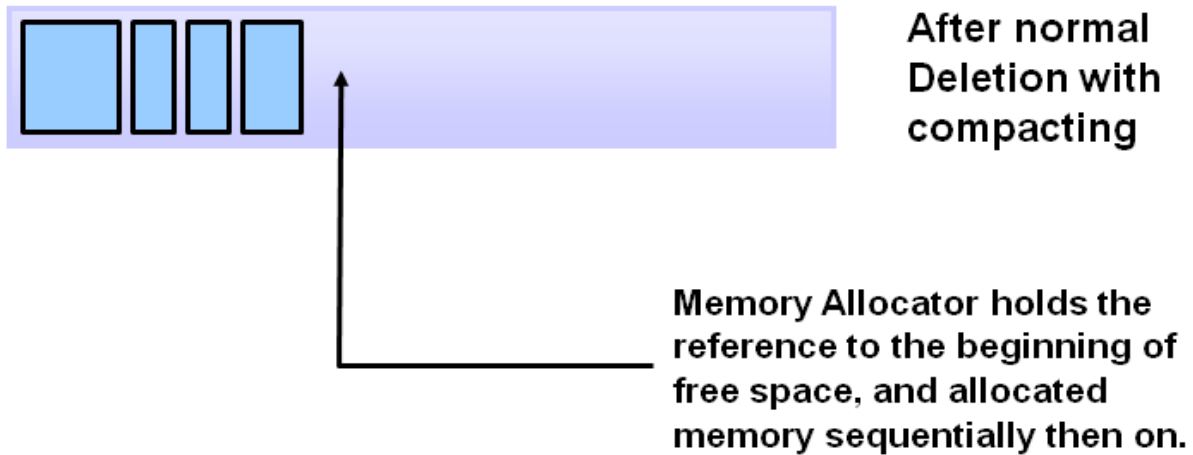


After normal deletion

Memory Allocator holds a list of references to free spaces, and searches for free space whenever an allocation is required

Mark and Sweep Garbage Collection

Deletion with Compacting



Outline

Memory Management in Java



Demo of Memory Analysis

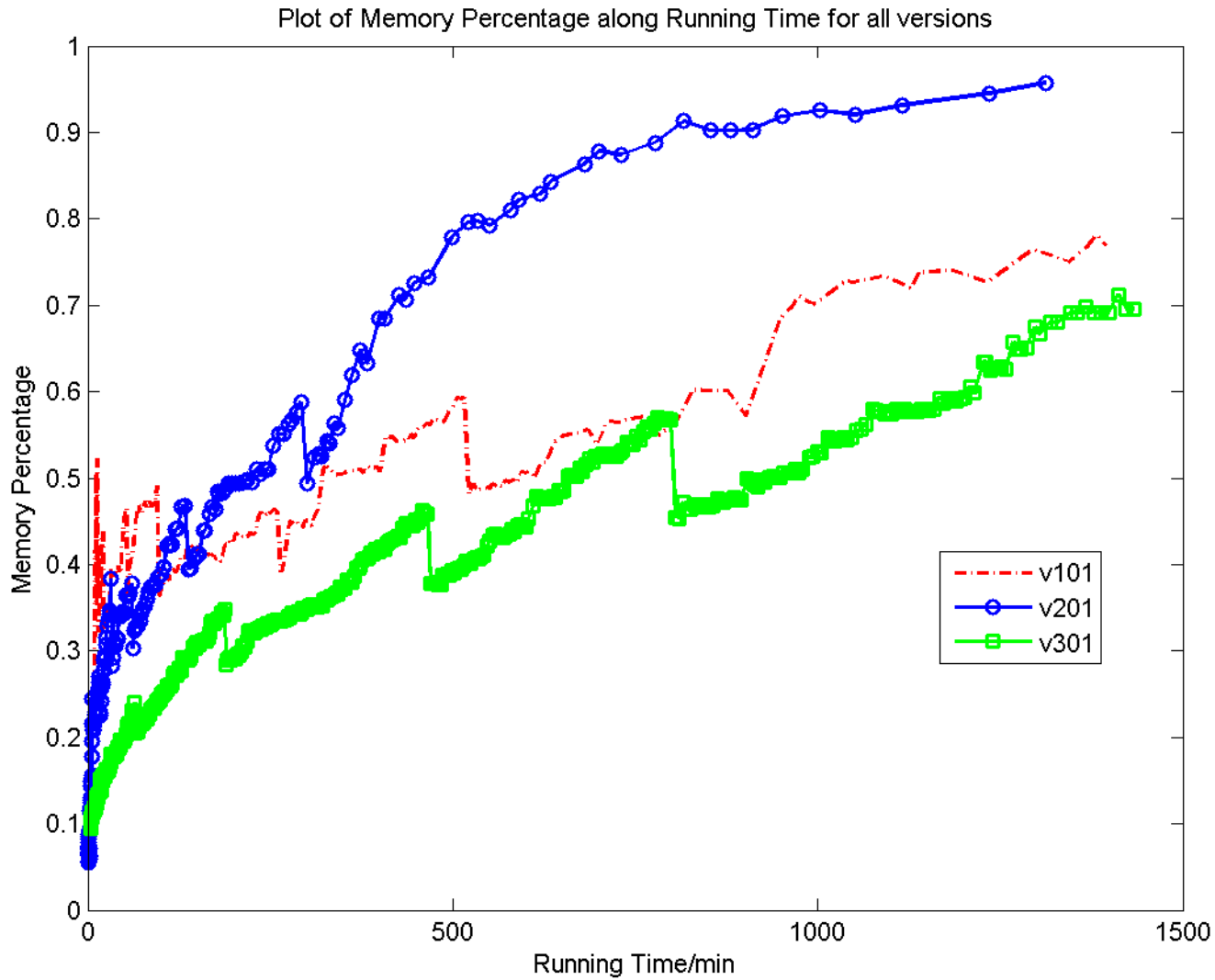


CPU-time Profiling

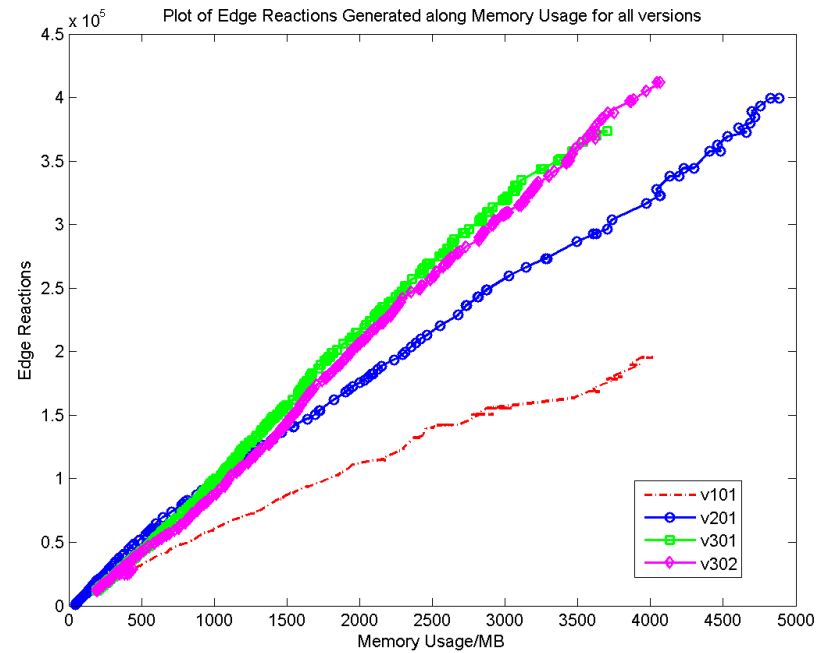
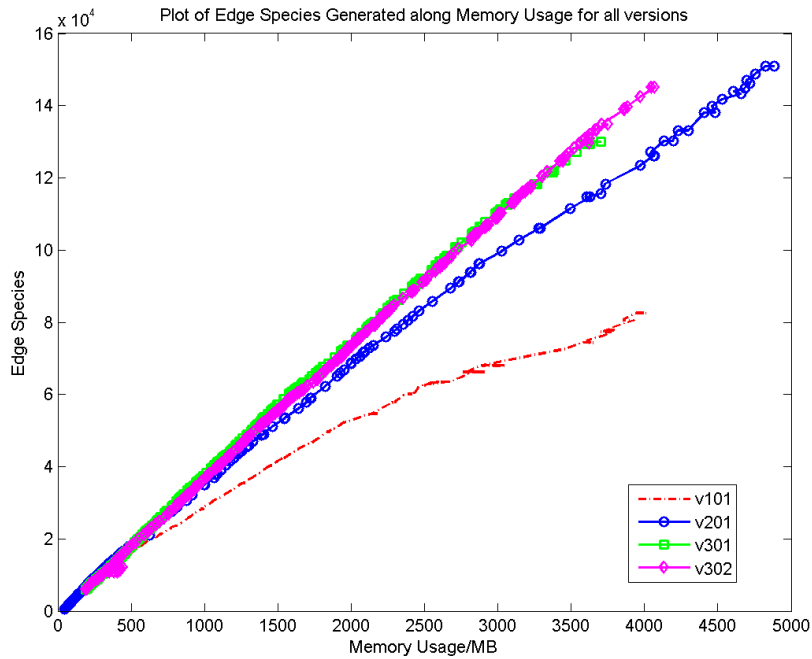


Possible Approaches

RAM limitation



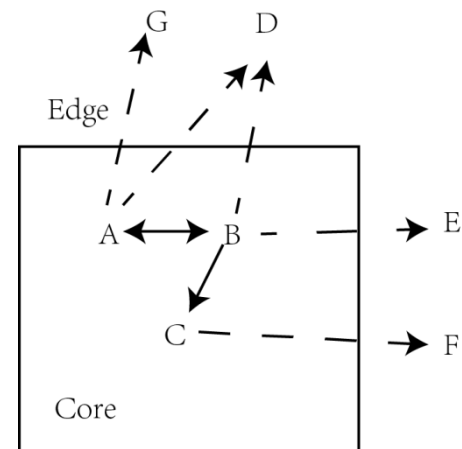
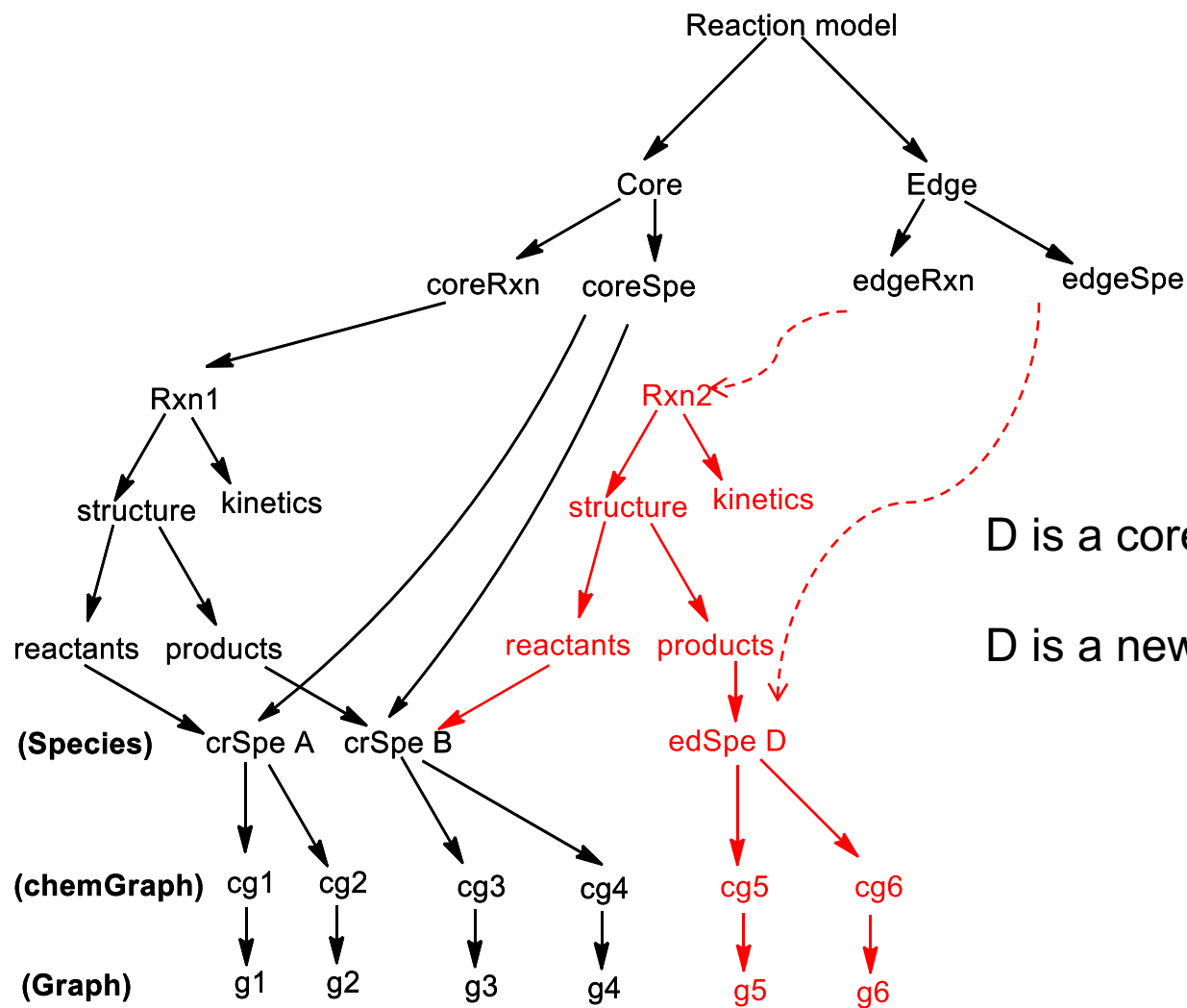
RAM limitation



Demo of Memory Analysis

- How to get a heap dump
 - Console: `jmap -dump:format=b,file=<filename.hprof> <pid>`
 - .sh file: `-XX:+HeapDumpOnOutOfMemoryError`
- How to use Eclipse Memory Analyzer
 - Histogram
 - Outgoing & incoming
 - Dominator tree & immediate dominator
 - Retained set

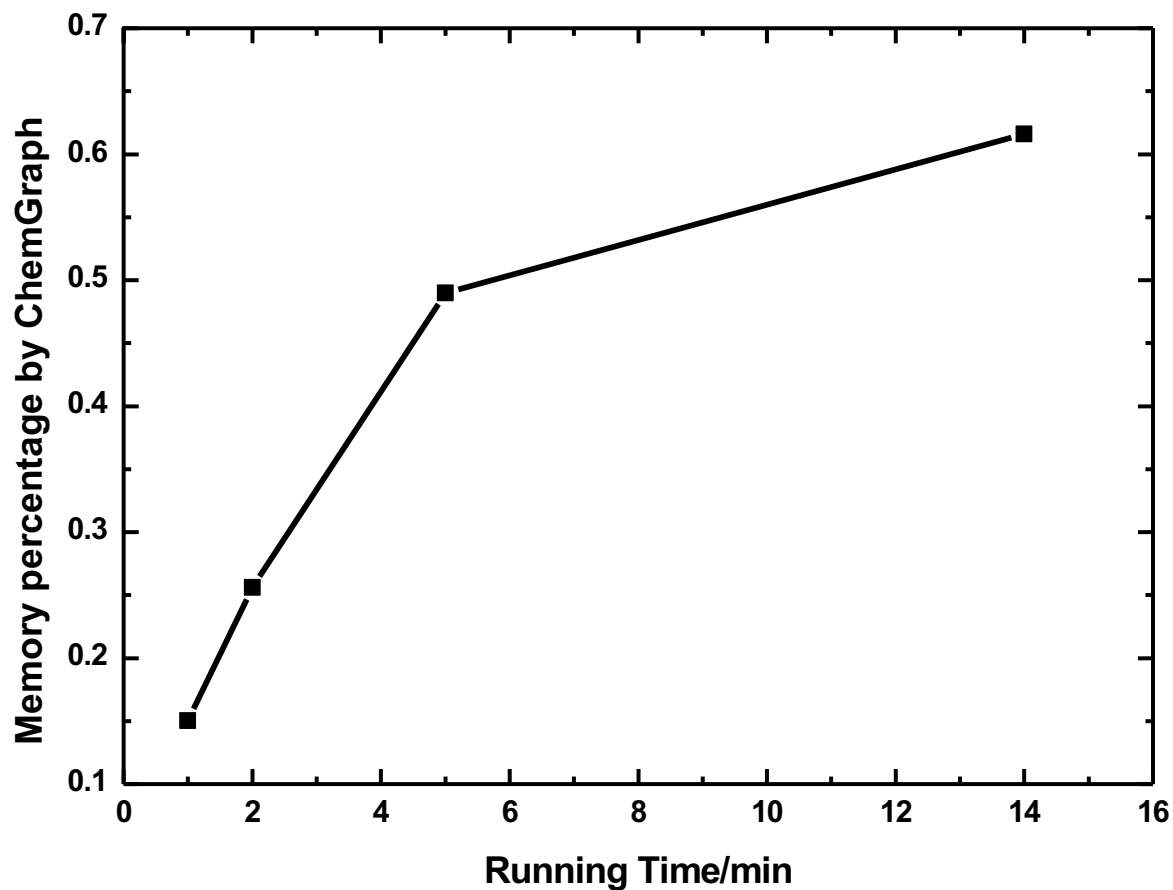
Object Graph in RMG-Java



D is a core species? **X**

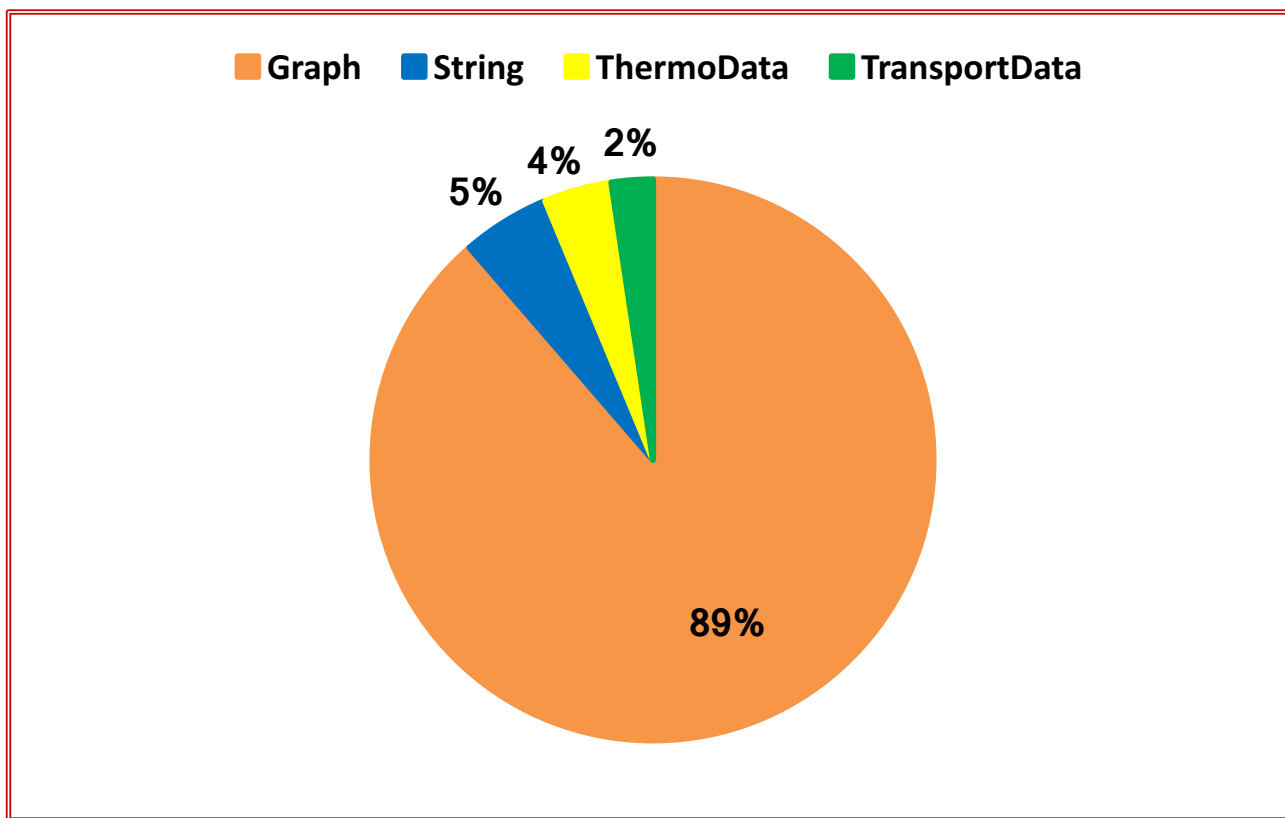
D is a new edge species? **✓**

Demo of Memory Analysis



ChemGraph is the class of objects that occupy most RAM!

What ChemGraph Dominates?



Outline

Memory Management in Java



Demo of Memory Analysis



CPU-time Profiling



Possible Approaches

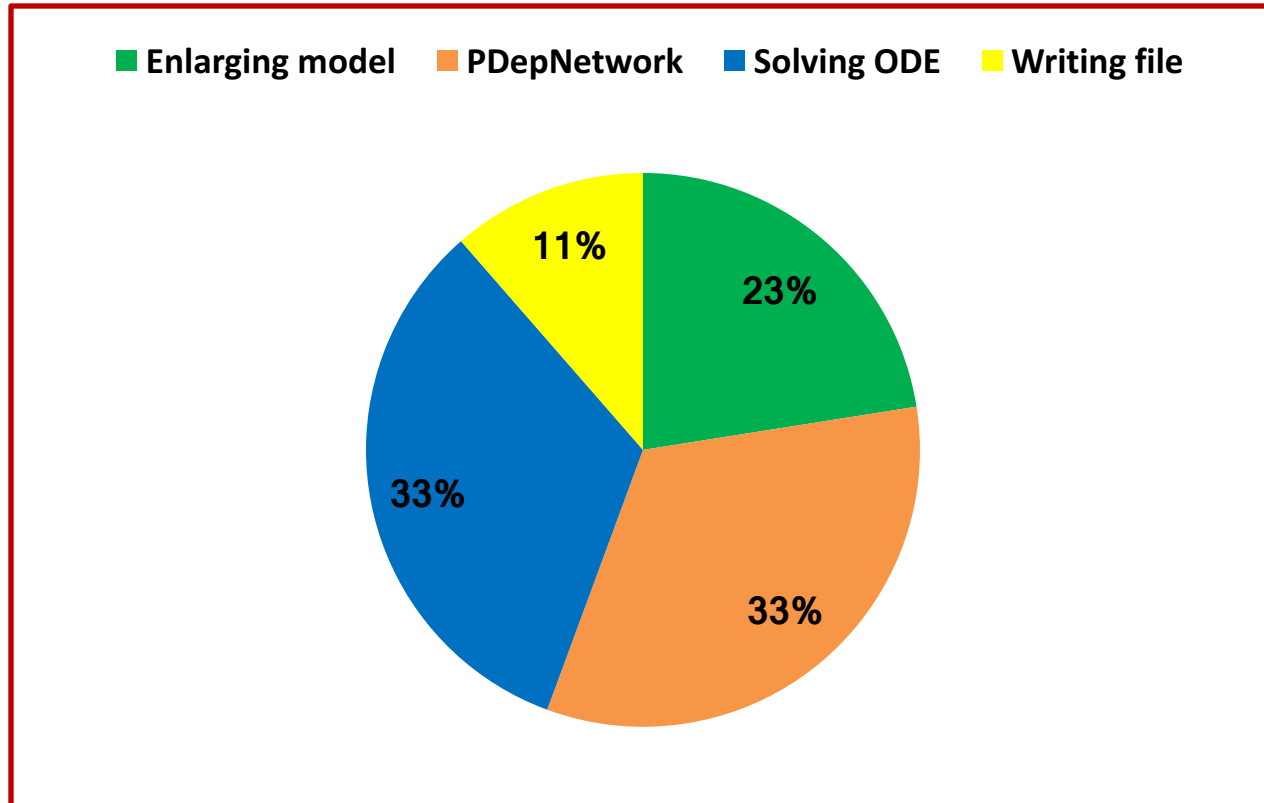
CPU-time Profiling

```
// ENLARGE THE MODEL!!! (this is where the good stuff happens)
pt = System.currentTimeMillis();
enlargeReactionModel();
double totalEnlarger = (System.currentTimeMillis() - pt) / 1000 / 60;
```

```

//*****333*****%% initialize PDepNetwork
pt = System.currentTimeMillis();
// 10/24/07 gmagoon: changed to use reactionSystemList
if ((reactionModelEnlarger instanceof RateBasedPDepRME)) { // 1/2/09 gmagoon and rwest: only call
PDepNetwork for P-dep cases
    for (Iterator iter = reactionSystemList.iterator(); iter
        .hasNext();) {
        ReactionSystem rs = (ReactionSystem) iter.next();
        rs.initializePDepNetwork();
    }
    // reactionSystem.initializePDepNetwork();
}
double initPDep = (System.currentTimeMillis() - pt) / 1000 / 60;
```

CPU-time Profiling



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CPU-time Profiling



Possible Approaches

Approach1: Memory Usage Reduction

At later stage of reaction generation:

- ChemGraph takes up **most** memory,
- **> 95%** ChemGraphs are for edge species.

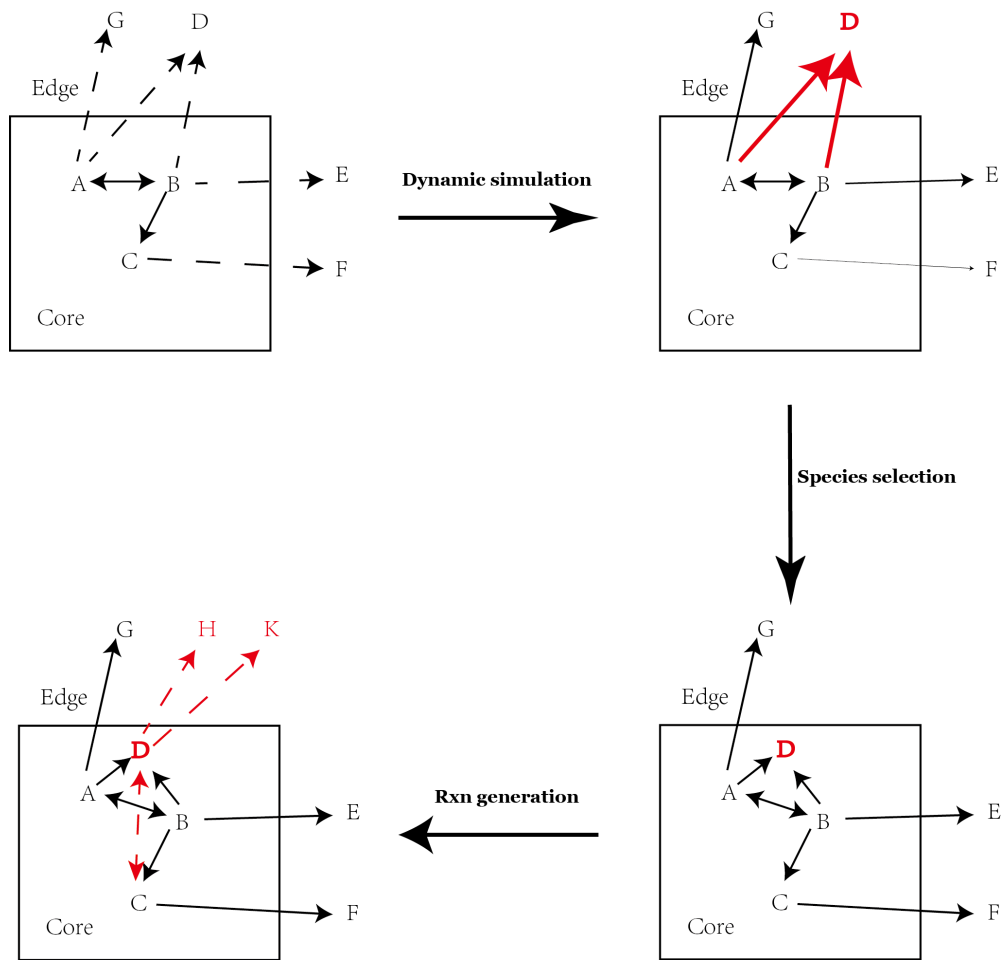
Most ChemGraphs occupy memory but contribute little

Proposed approach:

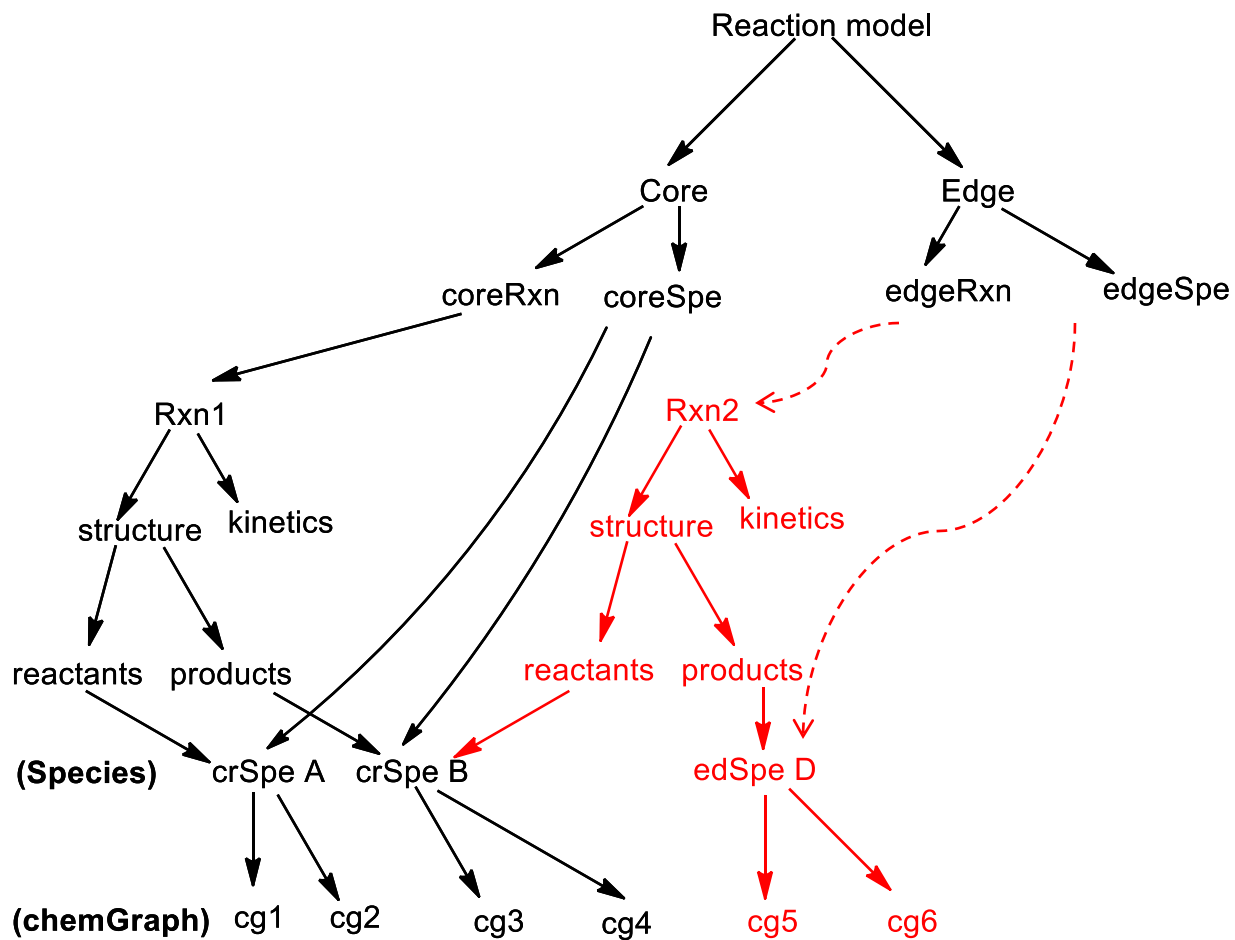
Replace edge's ChemGraphs with much cheaper identifiers

- One identifier < 100bytes, while one ChemGraph ~ 10^4 bytes,
- Can retrieve ChemGraphs back when needed,
- Can compare with other edge species using identifiers.

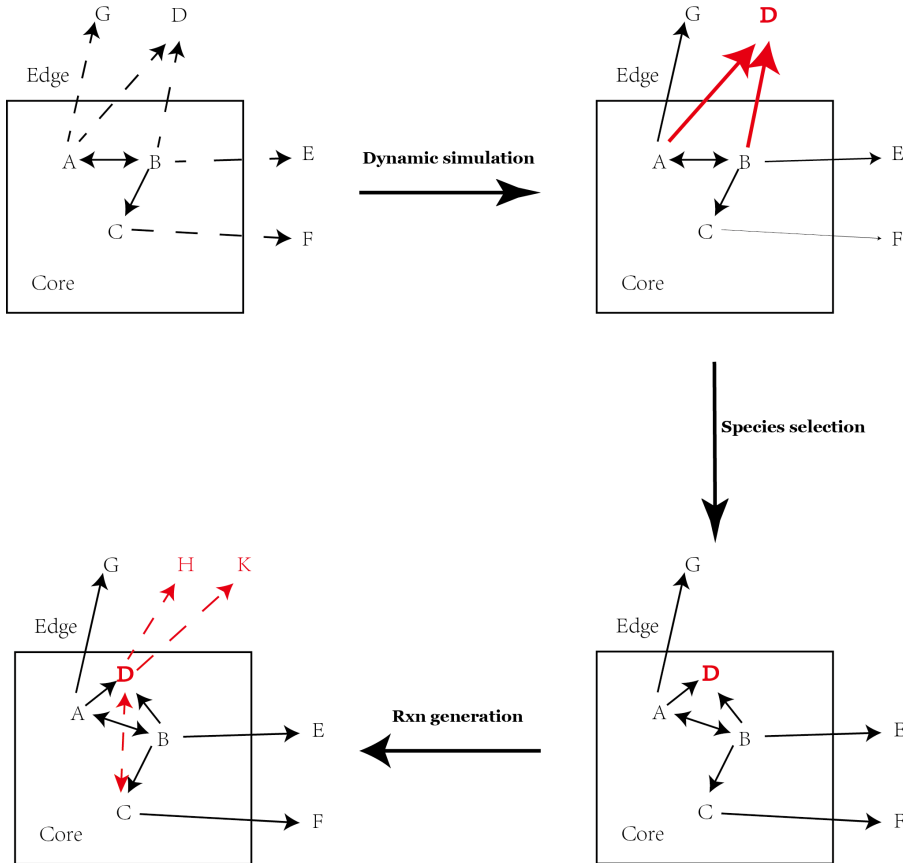
One iteration from view of MEMORY



Upon Reaction Generation



New steps added

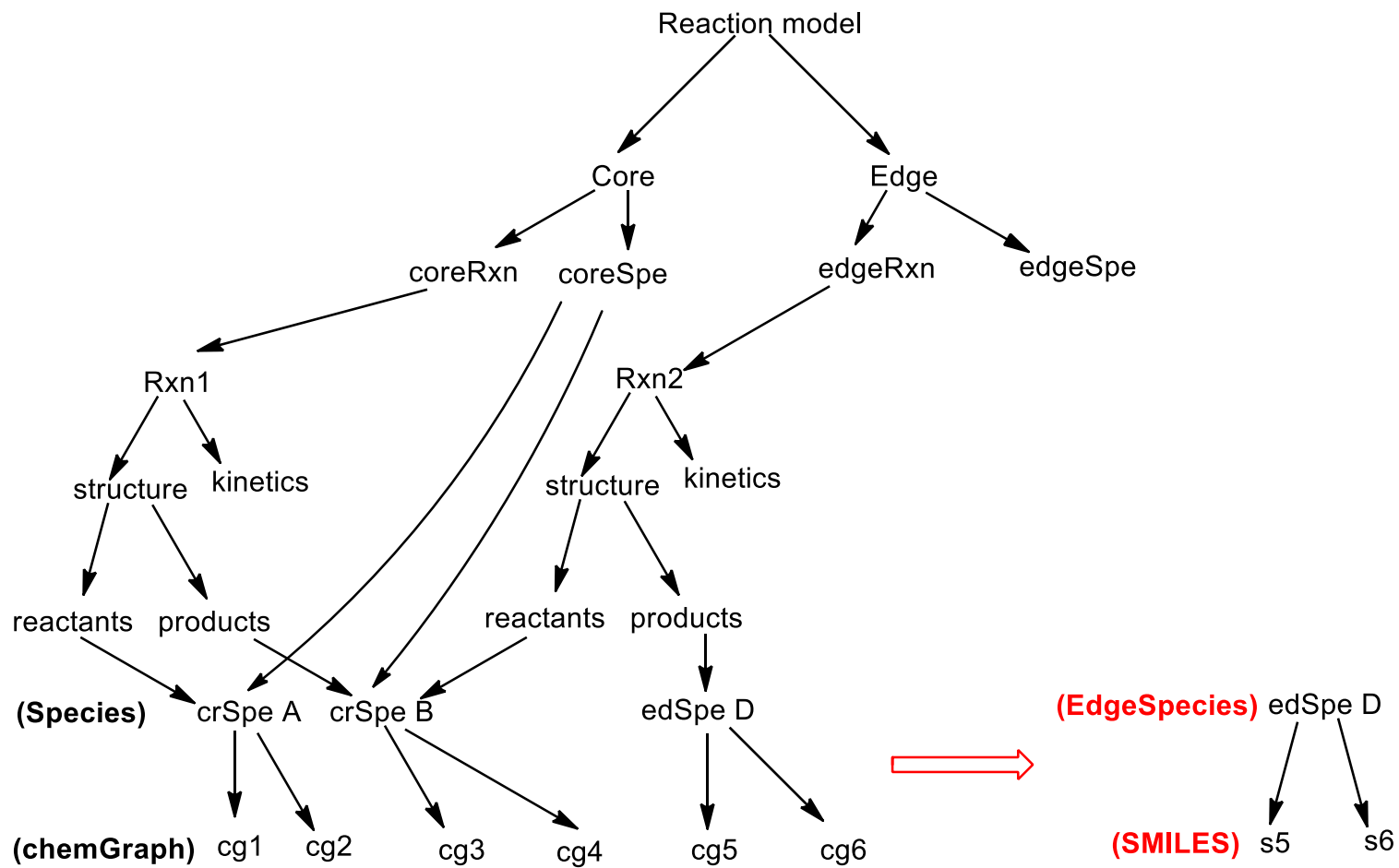


In original design, **Dynamic simulation** is the next step;

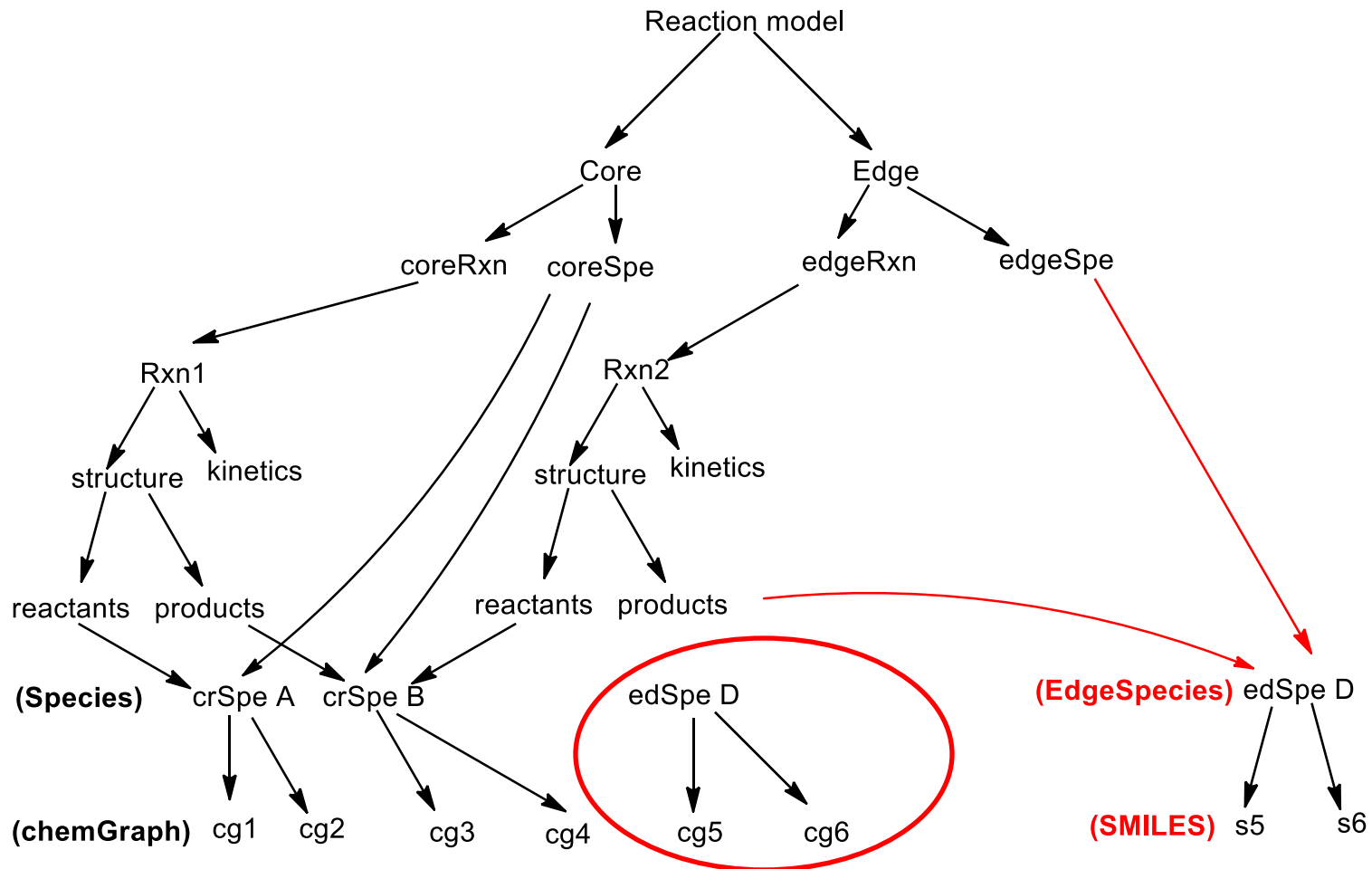
Now new steps added *BEFORE* that:

Memory Usage Reduction Method

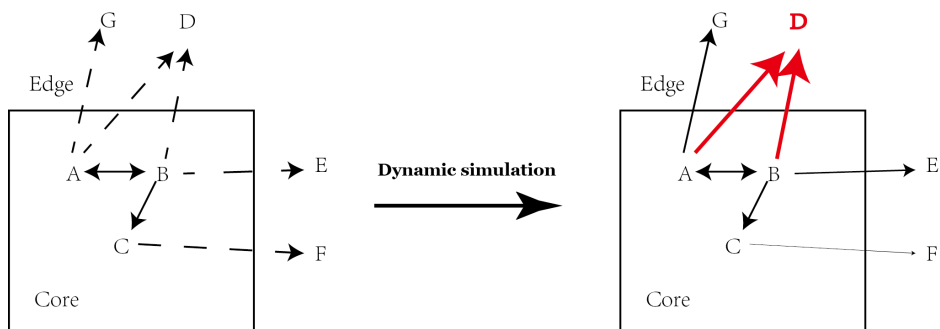
ChemGraph → SMILES



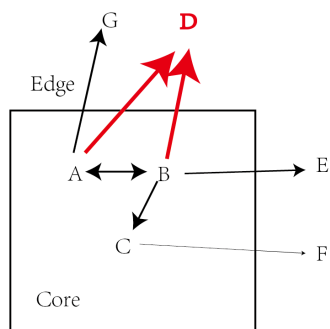
If edge species D is a new one



Garbage collected!



Dynamic simulation

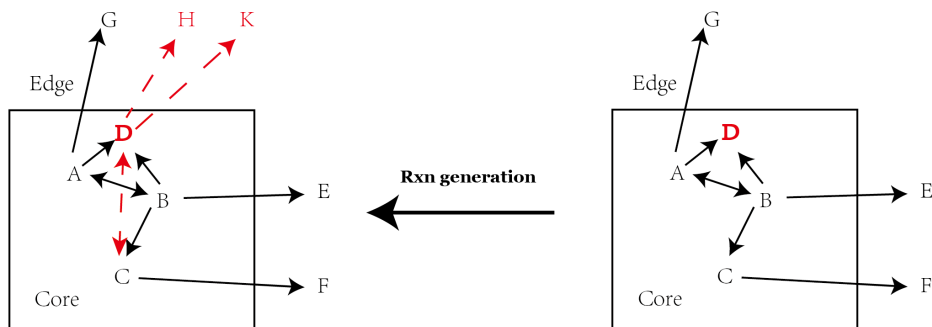


Now comes

Dynamic simulation & Selection

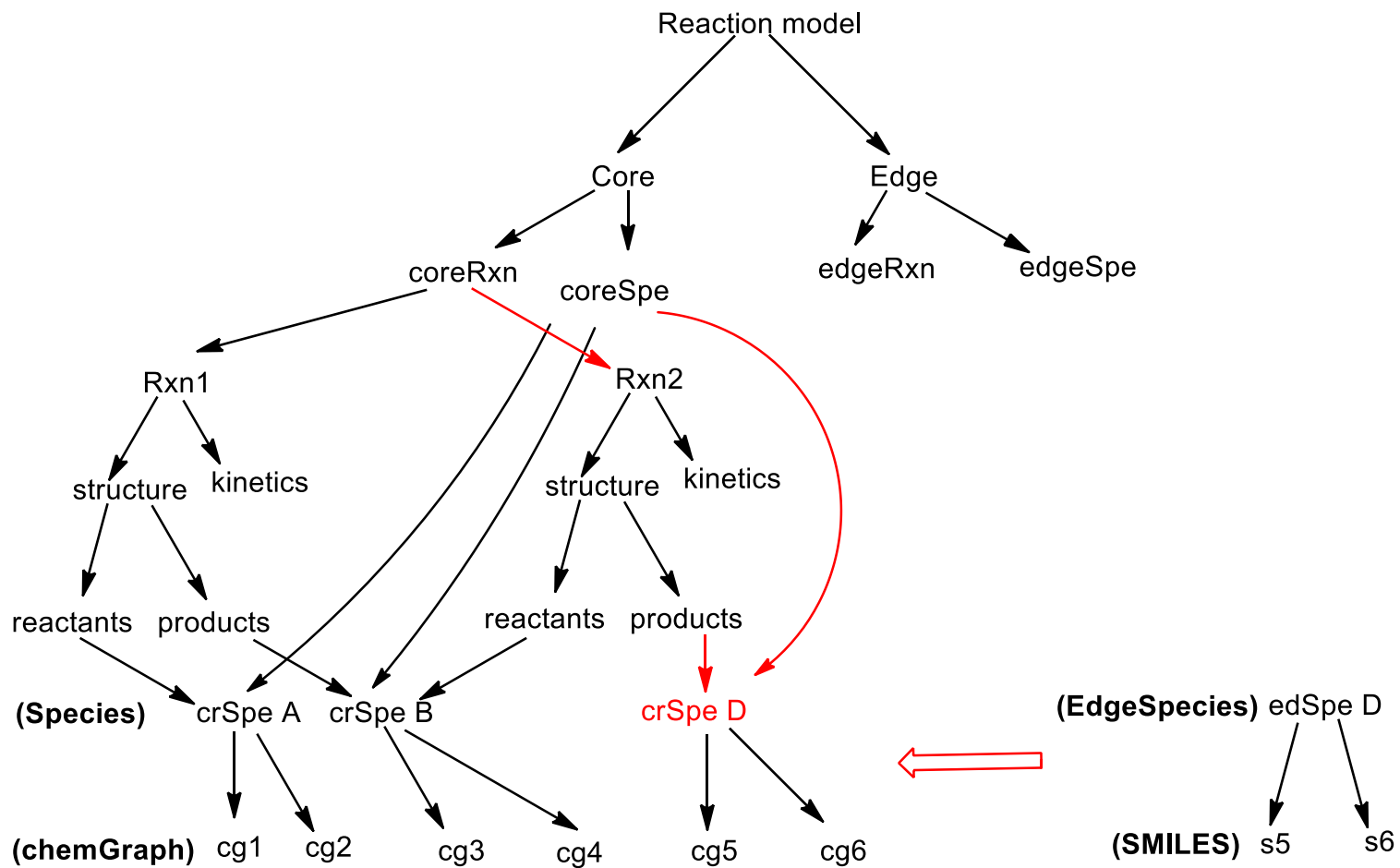
Species selection

Edge species D will be finally entering core

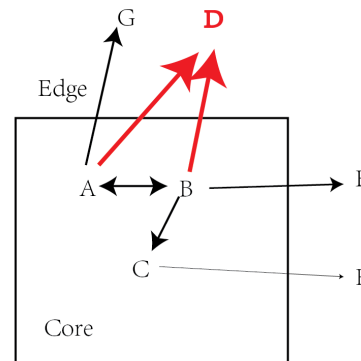
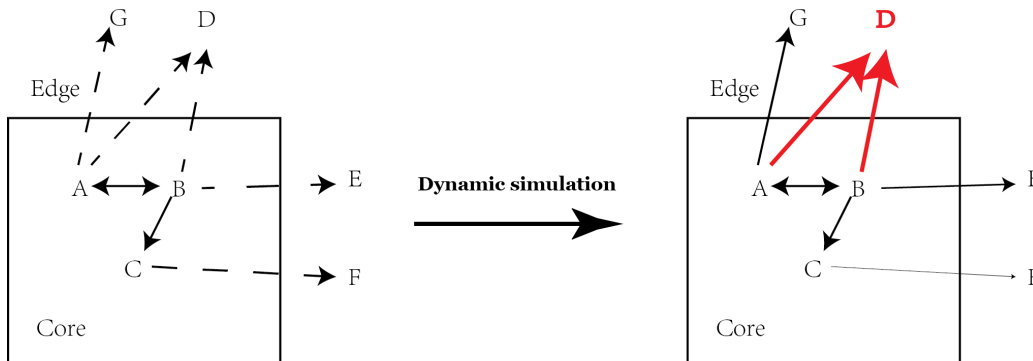


Rxn generation

Upon Species D being selected



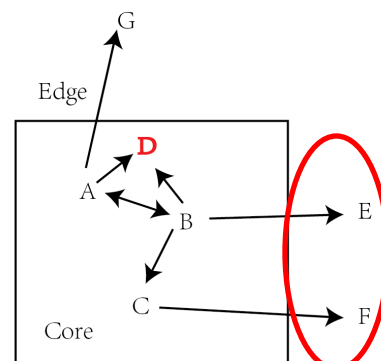
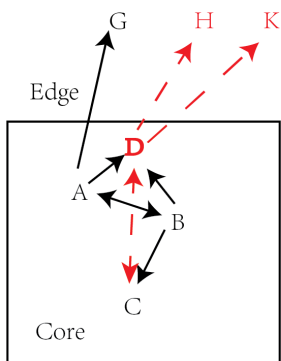
Approach2: Pruning Edge Species



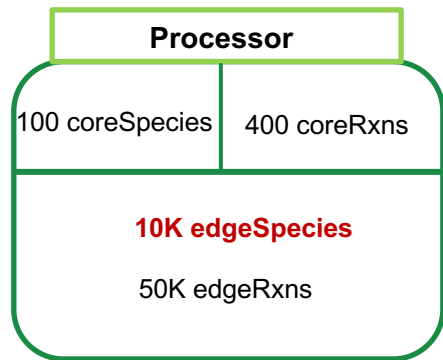
Pruning will be done based on fluxes.

- Upper limit of edge species
- Below a certain flux

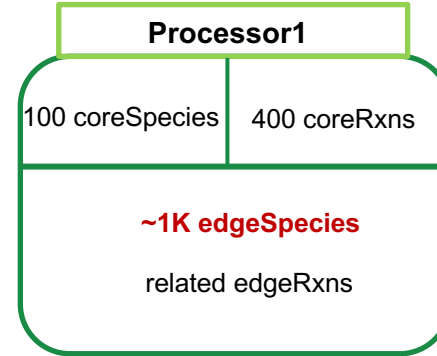
Species selection



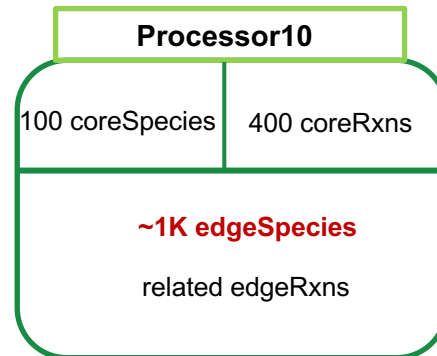
Approach3: Job Partition



spread to 10 processors



.....



Heavily limited by the 10K edge species

How to Partition Job

- Each processor **keeps a copy** of core model in its own memory;
- Edge species **almost evenly split** into N pieces for N processors;
 - Using M.W. makes partition easy and fast
 - Processor1 collects those species with $M.W. \leq 30$
 - Processor2 collects those with $30 < M.W. \leq 60$
 -
- Edge reactions **go where corresponding** edge species **go**;
 - e.g. $CH_3 + C_2H_6 \rightarrow CH_4 (M.W.=16) + C_2H_5 (M.W.=29)$
should go to Processor1
 - e.g. $CH_3 + C_2H_5OH \rightarrow CH_4 (M.W.=16) + C_2H_5O (M.W.=45)$
Processor1 stores $CH_3 + C_2H_5OH \rightarrow CH_4 (M.W.=16) + \text{“other edgeSpecies”}$
Processor2 stores $CH_3 + C_2H_5OH \rightarrow C_2H_5O (M.W.=45) + \text{“other edgeSpecies”}$

How Job Runs Differently

Step1: ODE solving. (Not affected)

- Edge species **don't serve as reactants**
- Core species and edge species are **decoupled** in ODE system
- ODE solver in each processor stops **at different conversion**

Step2: select new core species. (Need communication)

- Processor with **smallest conversion**

Step3: update core and edge model.

- Move the new core species from edge to core
- Move related edge reactions to core except those having “other edgeSpecies”
- Make reactions between new core species and old core species
 - Not all products are core species → checking where to go
 - All products are core species → checking reverse reactions

How Job Runs Differently

